## Graph Search Methods

- A vertex $u$ is reachable from vertex $v$ iff there is a path from v to u .



## Graph Search Methods

- A search method starts at a given vertex v and visits/labels/marks every vertex that is reachable from v .



## Graph Search Methods

- Many graph problems solved using a search method.
- Path from one vertex to another.
- Is the graph connected?
- Find a spanning tree.
- Etc.
- Commonly used search methods:
- Depth-first search.
- Breadth-first search.


## Depth-First Search

DFS(v)<br>\{<br>Label vertex v as reached.<br>for (each unreached vertex $u$ adjacenct from v)

DFS(u);
\}

## Depth-First Search Example



Start search at vertex 1.
Label vertex 1 and do a depth first search
from either 2 or 4.
Suppose that vertex 2 is selected.


Label vertex 2 and do a depth first search from either 3,5 , or 6 .

Suppose that vertex 5 is selected.

## Depth-First Search Example



Label vertex 5 and do a depth first search from either 3,7 , or 9 .
Suppose that vertex 9 is selected.


Label vertex 9 and do a depth first search from either 6 or 8 .
Suppose that vertex 8 is selected.

## Depth-First Search Example



Label vertex 8 and return to vertex 9 .
From vertex 9 do a DFS(6).


Label vertex 6 and do a depth first search from either 4 or 7 .

Suppose that vertex 4 is selected.

## Depth-First Search Example



Label vertex 4 and return to 6 .
From vertex 6 do a DFS(7).


Label vertex 7 and return to 6 .
Return to 9.

## Depth-First Search Example



Return to 5.

Depth-First Search Example


Do a DFS(3).

## Depth-First Search Example



Label 3 and return to 5 .
Return to 2.


Return to 1.

Depth-First Search Example


Return to invoking method.

In Class Exercise
Do Depth first search on the following tree


## Depth-First Search Property

- All vertices reachable from the start vertex (including the start vertex) are visited.


## Path From Vertex v To Vertex u

- Start a depth-first search at vertex v.
- Terminate when vertex $u$ is visited or when DFS ends (whichever occurs first).
- Time
- $\mathrm{O}\left(\mathrm{n}^{2}\right)$ when adjacency matrix used
- $\mathrm{O}(\mathrm{n}+\mathrm{e}$ ) when adjacency lists used (e is number of edges)


## Is The Graph Connected?

- Start a depth-first search at any vertex of the graph.
- Graph is connected iff all n vertices get visited.
- Time
- $\mathrm{O}\left(\mathrm{n}^{2}\right)$ when adjacency matrix used
- $\mathrm{O}(\mathrm{n}+\mathrm{e}$ ) when adjacency lists used (e is number of edges)


## Connected Components

- Start a depth-first search at any as yet unvisited vertex of the graph.
- Newly visited vertices (plus edges between them) define a component.
- Repeat until all vertices are visited.


## Connected Components



## Time Complexity

- $\mathrm{O}\left(\mathrm{n}^{2}\right)$ when adjacency matrix used
- $\mathrm{O}(\mathrm{n}+\mathrm{e})$ when adjacency lists used (e is number of edges)


## Spanning Tree



Depth-first search from vertex 1.
Depth-first spanning tree.

## Spanning Tree

- Start a depth-first search at any vertex of the graph.
- If graph is connected, the n-1 edges used to get to unvisited vertices define a spanning tree (depth-first spanning tree).
- Time
- $\mathrm{O}\left(\mathrm{n}^{2}\right)$ when adjacency matrix used
- $\mathrm{O}(\mathrm{n}+\mathrm{e}$ ) when adjacency lists used (e is number of edges)


## Breadth-First Search

- Visit start vertex and put into a FIFO queue.
- Repeatedly remove a vertex from the queue, visit its unvisited adjacent vertices, put newly visited vertices into the queue.


## Breadth-First Search Example



Start search at vertex 1.

Breadth-First Search Example


Visit/mark/label start vertex and put in a FIFO queue.

Breadth-First Search Example


Remove 1 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



FIFO Queue 24

Remove 1 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 2 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 2 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 4 from Q; visit adjacent unvisited vertices; put in Q .

Breadth-First Search Example


Remove 4 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 5 from Q; visit adjacent unvisited vertices; put in Q .

Breadth-First Search Example


Remove 5 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 3 from Q; visit adjacent unvisited vertices; put in Q .

Breadth-First Search Example


Remove 3 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 6 from Q; visit adjacent unvisited vertices; put in Q .

Breadth-First Search Example


Remove 6 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 9 from Q; visit adjacent unvisited vertices; put in Q .

Breadth-First Search Example


FIFO Queue 78

Remove 9 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 7 from Q; visit adjacent unvisited vertices; put in Q .

Breadth-First Search Example


FIFO Queue 8

Remove 7 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



Remove 8 from Q; visit adjacent unvisited vertices; put in Q .

## Breadth-First Search Example



FIFO Queue

Queue is empty. Search terminates.

## Time Complexity

- Each visited vertex is put on (and so removed from) the queue exactly once.
- When a vertex is removed from the queue, we examine its adjacent vertices.
- $\mathrm{O}(\mathrm{n})$ if adjacency matrix used
- O (vertex degree) if adjacency lists used
- Total time
- $\mathrm{O}(\mathrm{mn})$, where m is number of vertices in the component that is searched (adjacency matrix)


## Time Complexity

- $\mathrm{O}(\mathrm{n}+$ sum of component vertex degrees) (adj. lists)
$=\mathrm{O}(\mathrm{n}+$ number of edges in component $)$


## Breadth-First Search Properties

- Same complexity as DFS.
- Same properties with respect to path finding, connected components, and spanning trees.
- There are problems for which bfs is better than dfs and vice versa.


## Homework

- Sec. 6.2 Exercise2 @P352

