

- Linear list.
- One end is called front.
- Other end is called rear.
- Additions are done at the rear only.
- Removals are made from the front only.
- FIFO (First In First Out)

Bus Stop Queue



Bus Stop Queue



Bus Stop Queue







Revisit Of Stack Applications Bus Stop Queue • Applications in which the stack cannot be replaced with a queue. Parentheses matching. • Towers of Hanoi. front Method invocation and return. • Application in which the stack may be replaced with a queue. • Rat in a maze. • Results in finding shortest path to exit. Lee's Wire Router Wire Routing start pin end pin



Represent as a grid in which components and already placed wires are denoted by blocked grid positions. (Can be used to solve the rat in the maze.)



Label all reachable squares 1 unit from start.



Label all reachable unlabeled squares 2 units from start.

Lee's Wire Router



Label all reachable unlabeled squares 3 units from start.



Label all reachable unlabeled squares 4 units from start.

Lee's Wire Router



		4							
		3							
		2							
		1	2						
		2							
	4	3	4						
I		4							
I									

Label all reachable unlabeled squares 5 units from start.



Label all reachable unlabeled squares 6 units from start.



End pin reached. Traceback.

Lee's Wire Router



	6	5	6							
		4	5							
3		3								
2		2								
1		1	2							
2		2		6						
3	4	3	4	5	6					
4		4	5	6						
5	6	5	6							
6										

End pin reached. Traceback.

Queue Operations

- IsEmpty ... return true iff queue is empty
- Front ... return front element of queue
- Rear ... return rear element of queue
- Push ... add an element at the rear of the queue
- Pop ... delete the front element of the queue

Queue in an Array

- Use a 1D array to represent a queue.
- Suppose queue elements are stored with the front element in queue[0], the next in queue[1], and so on.

Derive From arrayList

0 1 2 3 4 5 6

e

c d

- Pop() => delete queue[0], shift other elements one step left
 - O(queue size) time
- Push(x) => if there is capacity, add at right end
 -O(1) time

O(1) Pop and Push

• to perform each opertion in O(1) time (excluding array doubling), we use a circular representation.

Custom Array Queue

- Use a 1D array queue.
 - queue[]
- Circular view of array.



Custom Array Queue

• Possible configuration with 3 elements.



Custom Array Queue

• Another possible configuration with 3 elements.



Custom Array Queue

- Use integer variables front and rear.
 - front is one position counterclockwise from first element
 - rear gives position of last element



Push An Element

• Move rear one clockwise.



Push An Element

- Move rear one clockwise.
- Then put into queue[rear].



Pop An Element

• Move front one clockwise.



Pop An Element

- Move front one clockwise.
- Then extract from queue[front].



Moving rear Clockwise

• rear++;

if (rear = capacity) rear = 0;



• rear = (rear + 1) % capacity;



Empty That Queue



Empty That Queue



- When a series of removes causes the queue to become empty, front = rear.
- When a queue is constructed, it is empty.
- So initialize front = rear = 0.



A Full Tank Please



A Full Tank Please



- When a series of adds causes the queue to become full, front = rear.
- So we cannot distinguish between a full queue and an empty queue!

Ouch!!!!!

- Remedies.
 - Don't let the queue get full.
 - When the addition of an element will cause the queue to be full, increase array size.
 - This is what the text does.
 - Define a boolean variable lastOperationIsPush.
 - Following each push set this variable to true.
 - Following each pop set to false.
 - Queue is empty iff (front == rear) && !lastOperationIsPush
 - Queue is full iff (front == rear) && lastOperationIsPush

Ouch!!!!!

- Remedies (continued).
 - Define an integer variable size.
 - Following each push do size++.
 - Following each pop do size--.
 - Queue is empty iff (size == 0)
 - Queue is full iff (size == arrayLength)
 - Performance is slightly better when first strategy is used.

Doubling Queue Capacity [3] [2] rear Е D front С Before enlarge array [1] [4] 0 2 3 4 5 1 Α B [0] [5] В С Е D Α After enlarge array

0	1	2	3	4	5	6	7	8	9	10	11		
В	С	D	E		A								
Shift													
0	1	2	3	4	5	6	7	8	9	10	11		
В	С	D	E								Â		

Homework

- Sec. 3.5 Exercise 1 (a) P157
 - Trace the program. (Find a path through the maze with Lee's Wire Router algorithm introduced in this section)