## Queues

## Bus Stop Queue

- Linear list.
- One end is called front.
- Other end is called rear.
- Additions are done at the rear only.
- Removals are made from the front only.
- FIFO (First In First Out)



## Bus Stop Queue

Bus Stop Queue


## Bus Stop Queue



## Revisit Of Stack Applications

- Applications in which the stack cannot be replaced with a queue.
- Parentheses matching.
- Towers of Hanoi.
- Method invocation and return.
- Application in which the stack may be replaced with a queue.
- Rat in a maze.
- Results in finding shortest path to exit.


## Wire Routing



Represent as a grid in which components and already placed wires are denoted by blocked grid positions. (Can be used to solve the rat in the maze.)

## Lee’s Wire Router

start pinend pin


Label all reachable squares 1 unit from start.

Lee's Wire Routerstart pinend pin

Store the
position in
the queue.

Label all reachable unlabeled squares 2 units from start.

Lee’s Wire Routerstart pinend pin

Label all reachable unlabeled squares 3 units from start.


Lee’s Wire Routerstart pinend pin

Lee’s Wire Router
start pinend pin

Label all reachable unlabeled squares 5 units from start.

## Lee’s Wire Router

Lee's Wire Routerstart pinend pin

start pin
end pin


End pin reached. Traceback.

Lee’s Wire Routerstart pinend pin


## Queue Operations

- IsEmpty ... return true iff queue is empty
- Front ... return front element of queue
- Rear ... return rear element of queue
- Push ... add an element at the rear of the queue
- Pop ... delete the front element of the queue

End pin reached. Traceback.

## Queue in an Array

- Use a 1D array to represent a queue.
- Suppose queue elements are stored with the front element in queue[0], the next in queue[1], and so on.


## O(1) Pop and Push

- to perform each opertion in $\mathrm{O}(1)$ time (excluding array doubling), we use a circular representation.


## Derive From arrayList

```
a a b b 
    llllllll
```

- $\operatorname{Pop}()=>$ delete queue[0], shift other elements one step left
- O(queue size) time
- Push(x) => if there is capacity, add at right end - O(1) time


## Custom Array Queue

- Use a 1D array queue.
queue[] $\square$
- Circular view of array.



## Custom Array Queue

- Possible configuration with 3 elements.



## Custom Array Queue

- Another possible configuration with 3 elements.



## Custom Array Queue

- Use integer variables front and rear.
- front is one position counterclockwise from first element
- rear gives position of last element



## Push An Element

- Move rear one clockwise.


## Push An Element

- Move rear one clockwise.
- Then put into queue[rear].


Pop An Element

- Move front one clockwise.
- Then extract from queue[front].



## Pop An Element

- Move front one clockwise.


Moving rear Clockwise

- rear++;
if (rear $==$ capacity) rear $=0$;

- rear = (rear + 1 ) \% capacity;


## Empty That Queue



## Empty That Queue



## Empty That Queue



## Empty That Queue



- When a series of removes causes the queue to become empty, front = rear.
- When a queue is constructed, it is empty.
- So initialize front $=$ rear $=0$.

A Full Tank Please


A Full Tank Please


A Full Tank Please


A Full Tank Please


- When a series of adds causes the queue to become full, front = rear.
- So we cannot distinguish between a full queue and an empty queue!


## Ouch!!!!!

- Remedies.
- Don’t let the queue get full.
- When the addition of an element will cause the queue to be full, increase array size.
- This is what the text does.
- Define a boolean variable lastOperationIsPush.
- Following each push set this variable to true.
- Following each pop set to false.
- Queue is empty iff (front == rear) \&\& !lastOperationIsPush
- Queue is full iff (front == rear) $\& \&$ lastOperationIsPush


## Ouch!!!!!

- Remedies (continued).
- Define an integer variable size.
- Following each push do size++.
- Following each pop do size--.
- Queue is empty iff (size == 0)
- Queue is full iff (size == arrayLength)
- Performance is slightly better when first strategy is used.

Doubling Queue Capacity (3) Before enlarge array

After enlarge array

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| B | C | D | E |  | A |  |  |  |  |  |  |


| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| B | C | D | E |  |  |  |  |  |  |  | A |

## Homework

- Sec. 3.5 Exercise 1 (a) P157
- Trace the program. (Find a path through the maze with Lee's Wire Router algorithm introduced in this section)

